



**APPLICATION FOR UNITED STATES  
LETTERS PATENT**

**CONTENT PROTECTION FOR A DIGITAL RECORDING DEVICE**

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By Express Mail # EL489904836US

## **BACKGROUND OF THE INVENTION**

### **1. Field of the Invention**

The present invention relates to recording devices for digital data streams, and particularly to methods and systems for restricting distribution of copies recorded on such devices.

### **2. Description of the Related Art**

Digital recording technology makes it possible to easily reproduce exact copies of digital data streams such as audiovisual material, including movies, television programs, and audio recordings. Compression technologies, such as specified by the Motion Pictures Expert Group (MPEG), reduce the requirements for data storage space so that it is possible to store such recordings on a hard disk, such as the hard disk in a personal computer (PC). Once content is captured on a hard disk, it is simple to make large numbers of copies on writable compact disks (CD's) or DVD disks, for example, or to distribute the content to others over network connections, such as over the Internet.

Such copying capability poses the threat that illegal copying can become widespread, causing significant economic losses to the owners of proprietary rights to the content. There is thus a need to protect recorded content from illicit reproduction and distribution, i.e., copyright infringement.

### **SUMMARY OF THE INVENTION**

In accordance with the present invention, digitally recorded content is protected by encrypting the data prior to recording it on a hard disk using a standard encryption method with an encryption key that varies from one digital recorder unit to another. When the data is played back, it is decrypted with the same key. An attempt to play back the encrypted file on another digital player or recorder unit will produce no intelligible output.

Manufacturing a number of recorders, each with a unique key completely hard-wired, is labor intensive and thus expensive. In the present invention, the key is formed by combining a relatively small portion hard-wired in the recorder and a large portion stored in a memory associated with the recorder and retrievable under control of software.

In an alternative embodiment a sub-portion of the large portion is entered by a user. This enables the user to preclude other users of the same recorder who do not know the sub-portion from playing back certain program content.

An aspect of the invention provides a one-time-programmable device for storing the encryption key, a first security bit, and a second security bit. A readout path is provided to retrieve the encryption key for purposes of testing after manufacture. When testing is complete, the first security bit is programmed ON (an irrevocable procedure). The readout path is arranged to be disabled by the ON condition of the first security bit. The readout path is further arranged to be re-enabled by programming ON the second security bit (also an irrevocable procedure); however another path essential to functioning of the recorder is arranged to be disabled by the ON condition of the second security bit. Retrieval of the key to program it into a second

recorder to replace the first one if the first one has broken down is thus possible, but retrieval of the key for illicit purposes is thus discouraged.

Other objects and features of the present invention will become apparent from the following detailed description considered in conjunction with the accompanying drawings. It is to be understood, however, that the drawings are designed solely for purposes of illustration and not as a definition of the limits of the invention, for which reference should be made to the appended claims. It should be further understood that the drawings are not necessarily drawn to scale and that, unless otherwise indicated, they are merely intended to conceptually illustrate the structures and procedures described herein.

### **BRIEF DESCRIPTION OF THE DRAWINGS**

In the drawings, wherein like reference numerals denote similar elements throughout the several views:

Fig. 1 is a high-level block diagram of a system including a recorder embodying the  
5 present invention;

Fig 2 shows further detail of a content protection device shown in the recorder of Fig. 1;

Fig. 3 depicts a memory address allocation scheme employed to enhance features of the present invention;

Fig. 4 depicts a basic scheme for recovering an internal key from the content protection  
10 device shown in Fig. 2; and

Fig. 5 depicts a scheme according to the present invention for recovering an internal key from the content protection device shown in Fig. 2.

## **DETAILED DESCRIPTION OF THE PRESENTLY PREFERRED EMBODIMENTS**

Fig. 1 is a high-level block diagram of a system embodying the present invention.

Recorder 1 includes digital processor 100 which receives a digital data stream, which for purposes of the present example is a digital television broadcast, but which might be digitized audio, digitized video such as a movie or portion thereof, digitized information, etc.

Transmission of the digital data stream to digital processor 100 may be by means of radio-frequency transmission, cable, or the like. Digital processor 100 processes the digital data stream into a form intelligible to a user; in the present example a digital TV broadcast is presented on TV monitor 500.

The user of recorder 1 may direct that the digital data stream be recorded for later use. Such recording may be on video cassette or rewritable compact disc (CD/R). Use of compression techniques such as MPEG-2 (a standard established by the Motion Pictures Expert Group) enables recording on a hard disk drive. In a present embodiment, recorder 1 is associated with a personal computer (PC) 400, with which a hard drive 420 is associated; Recorder 1 communicates with PC 400 via PCI bus 410 (so named because it conforms to the Peripheral Component Interconnect bus standard) and through PC 400 recorder 1 communicates with hard drive 420. In an alternative embodiment, recorder 1 is a self-contained set-top box relative to TV monitor 500, in which case storage 420 and PC 400 are located in box housing recorder 1. Other embodiments may be contemplated that do not include a personal computer. Such other embodiments would still have a storage device, such as hard drive 420.

Once program content is recorded on the storage device, copies of the content could then be distributed easily in the manner of transferring any PC files, such as making and distributing

multiple copies on multiple CD's, or distributing copies over the Internet 430. Since such distribution would be detrimental to owners of copyright in the content, the present invention stores encrypted copies in a manner adapted to limit such distribution.

Recorder 1 includes content protection device (CPD) 200 interposed between digital  
5 processor 100 and storage device 420 to perform the encryption and decryption. In a preferred embodiment, CPD 200 is a digital integrated circuit (ASIC). In a preferred embodiment, CPD 200 encrypts using the 3DES (data encryption standard) algorithm, with a 168-bit key. CPD 200 of any particular recorder 1 uses a particular key different from that used in other recorders 1. Thus, a copy given to the owner of another recorder 1 would not play back intelligibly on the  
10 other recorder 1.

Manufacturing such ASICs with unique on-chip 168-bit keys can be difficult and expensive. A preferred embodiment employs an alternative solution of using an externally applied unique or pseudo-unique key portion that can be stored in non-volatile memory such as FLASH memory 300, and an internal key portion "on chip" inside the said CPD 200. In a  
15 preferred embodiment the internal key portion is in a one-time-programmable ROM within CPD 200. In an alternative embodiment the internal key portion is encoded in bonding wires bonded during manufacture of the ASIC embodying CPD 200. The actual key used when encrypting and decrypting data is constructed as a combination of the external key and the internal key. The external key can be changed from time to time by software, for instance after each movie  
20 recorded. Key management software, running in digital processor 100, can maintain a table of the keys used in the FLASH memory, so that only indexes to the table need be passed to the CPD 200 to be attached to the file that contains the recording. The external keys can be made unique

for each unit, for example, by using the serial number of the unit as a seed for a key generation algorithm. Pseudo-unique keys can be generated by using a random number as the seed. Noise received by a satellite, cable, or terrestrial front-end that is not tuned to a proper channel can be used as a source for the random numbers. Variations to the internal key can be made during  
5 packaging of the integrated circuit through, for example, the use of different patterns for bonding wires inside the device. This design makes it possible to manufacture small batches of devices with different internal keys for a lower cost, relative to manufacturing devices with the entire key being unique.

At user direction, a sub-portion of the external key portion may be provided through user  
10 interface 350. Using this feature, a user may restrict other users of the same recorder 1 from accessing certain program content. For example, a parent may provide a key sub-portion when recording material inappropriate for children, so that his minor children are then unable to view that material because they do not know the sub-portion of the key.

Fig. 2 is a block diagram showing internal detail of CPD 200. (CPD 200 includes  
15 software-controlled logic (not shown) for effecting intercommunication among CPD 200's various components.) Digital processor interface 202 maintains two-way communication between CPD 200 and digital processor 100, as well as with flash memory 300 and user interface 350. Dual-port random-access memory (DPRAM) 210 stores data being transferred into or out of CPD 200, and is heavily used to transfer software-generated control messages between digital  
20 processor 100 and PC 400. Interrupt logic 212 passes interrupt signals between digital processor 100 and PC 400. A major use of such interrupt signals is to alert PC 400 to the availability of encrypted data during recording, and to the need for more data retrieval during playback. Control



and status registers 208 store control information and transfer control information, such as key components, into and out of CPD 200. Internal key 204 represents the 8-bit portion of the 168-bit 3DES key used for encryption and decryption. Key logic 206 combines the 160-bit external key portion obtained from flash memory 300 (or alternatively, from flash memory 300 in conjunction with user interface 350) with the 8-bit internal key 204 to produce, according to a predetermined relationship, the 168-bit 3DES key.

A current portion of a digital data stream (such as program content of a digitized TV broadcast) sent from digital processor 100 for storage on hard drive 420 is input to encrypter input FIFO 214, from which the payload portion is input to 3DES encrypter 218, along with an encryption key from key logic 206. Packet headers bypass encrypter 218 under control of encryption control logic 224 and are passed by multiplexer 226 along with the encrypted payload portion of the data stream. The resultant stream is stored in encrypter output FIFO 230 for subsequent output through PCI interface 250 to PCI bus 410 for transfer to hard drive 420. Keeping the header portions unencrypted facilitates such functions as fast forward and fast rewind.

A data stream retrieved from hard drive 420 is passed over PCI bus 410 and through PCI interface 250 into decrypter input FIFO 232. It is input to 3DES decrypter 228 along with a decryption key from key logic 206. The encrypted payload portion, combined with unencrypted packet headers, are combined by multiplexer 220 under control of decryption control logic 222, and the resultant data stream is input to decrypter output FIFO 216 for transmission through digital processor interface 202 to digital processor 100 for further transmission, such as to TV monitor 500.

The encryption is performed on 188-byte MPEG-2 Transport Packets in such a fashion that the 4-byte packet headers are kept unencrypted, while the remaining 184 bytes of the packet (often referred to as the payload) are encrypted. This provides the possibility to demultiplex the encrypted content into video and audio packets, for instance, according to the packet ID (PID) in the header, without having to perform decryption. It will also be possible to scan the file back and forth (fast forward and fast rewind) without sending unnecessary information through the decrypter to the video and audio decoders in digital processor 100, as it is possible to find the boundaries of packetized elementary stream (PES) packets according to the Payload Unit Start Indicator bit in the header. Audio packets can be skipped altogether in these modes. The encryption device performs a check for the packet length, and alerts the software controlling CPD 200 by an interrupt if an attempt is made to transfer packets that deviate from the standard 188 bytes.

A feature of the preferred embodiment facilitates replacing packet headers during retrieval of a stored data stream. As shown in Fig. 3, the output of encrypter output FIFO 216 (serial bytes) is mapped to occupy a plurality of addresses in address space, and that plurality of addresses is contiguous with the addresses of DPRAM 210. A packet is written to encrypter output FIFO 216 as usual. The header is read out by software, and a replacement header is written to top locations of DPRAM 210. A DMA transfer is initiated, spanning from the replacement header in the DPRAM 210 through the address range representing the number of payload positions to be read out of the FIFO 216. Synchronizing the advancing of the payload stream from FIFO 216 with reading out of the designated address range causes the replacement

header and the remaining part of the packet to be transferred to its destination in one DMA operation.

The invention solves the problem of illegal copying and distribution of digital media content, while still allowing the user to make a recording for his or her own, local use by  
5 encrypting the recording with a key unique to the particular recording device, obviating playback on other devices which would have different keys.

The combination of using external and internal keys allows flexibility in key management. The external key can be changed by software from time to time, e.g., after each movie recorded. An external key alone would provide only a limited amount of security, because  
10 it would be possible for a person trying to circumvent the content protection scheme to observe the external key on the bus through which it is loaded to the encryption / decryption device, and then decrypt the content of the hard disk, for instance by running DES or 3DES decryption software. In order to discourage people from distributing the content in the encrypted form along with the external key(s), the internal key varies from one recorder 1 to another, making it hard to  
15 find another decryption device that can decrypt the content.

Keeping the transport packet headers unencrypted is an essential feature that allows easy implementation of trick modes such as fast forward and fast rewind. If the entire packets were encrypted it would be necessary to perform decryption for all the data that is skipped during playback, too.

20 The packet header and payload length security checks prevent abuse of the above described feature.

The packet header replacement feature, using a single DMA spanning from DPRAM to FIFO, is faster than copying the entire packet into memory, performing the replacement there, and then using DMA to transfer the packet with the replaced header to the audio decoder. This invention solves a problem that is specific to the audio decoder solution used, which dictates that DMA must be used to place data to the audio decoder's input buffer.

The encryption protection provided by the present invention, however, might work against the rightful owner of the copies in the event that his recorder 1 breaks down. If he obtains another recorder 1 it will have a different internal key 204, so he will not be able to play his recorded copies on it. A procedure for recovering the internal key 204 so that it can be programmed into another recorder 1 is possible, but is subject to abuse. Such a scheme is found in Fig. 4, which shows further detail of one-time programmable ROM 203 (introduced in Fig. 2) and the path from it to key logic 206. One-time programmable ROM 203 has at least eight bits, OTP1 through OTP8. Each of these eight bits is input to one of the eight AND gates 604, shown schematically as a single gate with the legend "x8". The gates are provided on the ASIC that comprises CPD 200. A signal directing that the internal key be passed to key logic 206 is provided at the other input of each of the eight AND gates 604, and the eight-bit key is passed through to key logic 206. As noted above in connection with Fig. 2, key logic 206 passes the key and/or derivative information to encrypter 218 and decrypter 228.

In a preferred embodiment, there is also provision for connecting an external IC tester 700. The tester may provide on line 710 a signal requesting the internal key, which is applied to an input of each of eight AND gates 602, each of which has its other input connected to a different bit of internal key 204. Internal key 204 is passed to the tester on lines 720, and the

tester may then determine that one-time programmable ROM 203 was programmed correctly as intended. The connections 710 and 720 could be used by a device other than tester 700 for requesting and obtaining internal key 204, and thus could be used by a device to retrieve internal key 204 which could then be programmed into another recorder 1 for use by a rightful owner whose original recorder 1 has broken down. However, that facility could also be used by a pirate wanting to gain illicit access to the copies on hard drive 420 for playback on a different recorder 1.

Fig. 5 shows another embodiment using two additional bits available from one-time programmable ROM 203, namely OTP9 and OTP10. The AND gates corresponding to 602 of Fig. 4 are here provided with an additional input each, and are named 602A. OR gate 606 is introduced in Fig. 5. During initial manufacture of a recorder 1, an internal key 204 is programmed into bits OTP1 through OTP8. OTP9 and OTP10 are not programmed to values of 1, and thus contain values of 0. The 0 state of OTP9 conditions OR gate 606 to enable an input of AND gates 602A, which can then function as described in connection with Fig. 4 to pass internal key 204 to tester 700 on lines 720.

When testing of CPD 200 is complete, bit OTP9 is programmed to a value of 1. OR gate 206 is then not conditioned, and AND gates 602A thus can not be enabled to pass the internal key out on lines 720. Anyone wishing to access the copies on hard drive 420 is thwarted. There is no effect on any other portion of recorder 1, so recorder 1 functions normally for its rightful owner.

In the event that recorder 1 breaks down and its owner wishes to obtain another one, it will be necessary to retrieve internal key 204 from his old unit in order to program it into one-

time programmable ROM 203 of his new unit, so that he will be able to play back copies that he had recorded using his old unit. A service facility programs bit OTP10 of the old recorder 1 to a value of 1, which conditions OR gate 606 to enable AND gates 602A, permitting output of the key on lines 720 so that it can be programmed into one-time programmable ROM 203 of the new recorder 1. However, under the present invention the ON or 1 state of OTP10 irrevocably disables the normal operation of the recorder 1. The actual manner of doing this is a design choice; in the example given in Fig. 5, the disabling is accomplished by disabling the recording and playback paths of CPD 200. As shown in Fig. 2, encrypter 218 and decrypter 228 are integral to the recording and playback paths respectively. As shown in Fig. 5, bit OTP10 is required to be in the OFF or 0 state to enable encrypter 218 and decrypter 228; programming bit OTP10 to the 1 state in order retrieve the key thus disables the operation of the recorder 1. Since programming bit OTP10 to the 1 state is an irrevocable operation, operation of recorder 1 is permanently disabled. Many alternative ways may be contemplated in which the 1 state of bit OTP10 may be deployed to disable the recorder 1.

Thus, while there have been shown and described and pointed out fundamental novel features of the invention as applied to a preferred embodiment thereof, it will be understood that various omissions and substitutions and changes in the form and details of the devices illustrated, and in their operation, may be made by those skilled in the art without departing from the spirit of the invention. For example, it is expressly intended that all combinations of those elements and/or method steps which perform substantially the same function in substantially the same way to achieve the same results are within the scope of the invention. Moreover, it should be recognized that structures and/or elements and/or method steps shown

and/or described in connection with any disclosed form or embodiment of the invention may be incorporated in any other disclosed or described or suggested form or embodiment as a general matter of design choice. It is the intention, therefore, to be limited only as indicated by the scope of the claims appended hereto.